

ω_1 under Π_1 -Collection

Toshiyasu Arai

Graduate School of Science, Chiba University
 1-33, Yayoi-cho, Inage-ku, Chiba, 263-8522, JAPAN
 tosarai@faculty.chiba-u.jp

Abstract

We describe a proof-theoretic bound on Σ_2 -definable countable ordinals in Kripke-Platek set theory with Π_1 -Collection and the existence of ω_1 .

1 Introduction

Let (ω_1) denote an axiom stating that ‘there exists an uncountable regular ordinal’, and $T_1 := \text{KP}\omega + (V = L) + (\Pi_1\text{-Collection}) + (\omega_1)$. Let ρ_0 denote the least ordinal above ω_1 such that $L_{\rho_0} \models (\Pi_1\text{-Collection})$. In this note a collapsing function $\Psi_{\omega_1} : \alpha \mapsto \Psi_{\omega_1}(\alpha) < \omega_1$ is defined, and it is shown that for each $n < \omega$, $T_1 \vdash \forall \alpha < \omega_n(\rho_0 + 1) \exists x < \omega_1(x = \Psi_{\omega_1}(\alpha))$ with a Σ_2 -formula $x = \Psi_{\omega_1}(\alpha)$, cf. Lemma 4.5. Conversely we show the

Theorem 1.1 *For a sentence $\exists x \in L_{\omega_1} \varphi(x)$ with a Σ_2 -formula $\varphi(x)$, if*

$$T_1 \vdash \exists x \in L_{\omega_1} \varphi(x)$$

then

$$\exists n < \omega [T_1 \vdash \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi(x)].$$

This paper relies on our [1].

2 Σ_1 -Skolem hulls

Everything in this section is reproduced from [1].

For a model $\langle M; \in \rangle$ and $X \subset M$, $\Sigma_1^M(X)$ denotes the set of $\Sigma_1(X)$ -definable subsets of M , where $\Sigma_1(X)$ -formulae may have parameters from X . $\Sigma_1^M(M)$ is denoted $\Sigma_1(M)$.

An ordinal $\alpha > 1$ is said to be a *multiplicative principal number* iff α is closed under ordinal multiplication, i.e., $\exists \beta[\alpha = \omega^{\omega^\beta}]$. If α is a multiplicative principal number, then α is closed under Gödel’s pairing function j and there exists a Δ_1 -bijection between α and L_α for the constructible hierarchy L_α up to α . In this section σ is assumed to be a multiplicative principal number $> \omega$.

Definition 2.1 1. $cf(\kappa) := \min\{\alpha \leq \kappa : \text{there is a cofinal map } f : \alpha \rightarrow \kappa\}.$

2. $\rho(L_\sigma)$ denotes the Σ_1 -projectum of L_σ : $\rho(L_\sigma)$ is the least ordinal ρ such that $\mathcal{P}(\rho) \cap \Sigma_1(L_\sigma) \not\subset L_\sigma$.
3. Let $\alpha \leq \beta$ and $f : L_\alpha \rightarrow L_\beta$. Then the map f is a Σ_1 -elementary embedding, denoted $f : L_\alpha \prec_{\Sigma_1} L_\beta$ iff for any $\Sigma_1(L_\alpha)$ -sentence $\varphi[\bar{a}] (\bar{a} \subset L_\alpha)$, $L_\alpha \models \varphi[\bar{a}] \Leftrightarrow L_\beta \models \varphi[f(\bar{a})]$ where $f(\bar{a}) = f(a_1), \dots, f(a_k)$ for $\bar{a} = a_1, \dots, a_k$. An ordinal γ such that $\forall \delta < \gamma [f(\delta) = \delta] \& f(\gamma) > \gamma$ is said to be the critical point of the Σ_1 -elementary embedding f if such an ordinal γ exists.
4. For $X \subset L_\sigma$, $\text{Hull}_{\Sigma_1}^\sigma(X)$ denotes the set (Σ_1 -Skolem hull of X in L_σ) defined as follows. $<_L$ denotes a Δ_1 -well ordering of the constructible universe L . Let $\{\varphi_i : i \in \omega\}$ denote an enumeration of Σ_1 -formulae in the language $\{\in\}$. Each is of the form $\varphi_i \equiv \exists y \theta_i(x, y; u)$ ($\theta \in \Delta_0$) with fixed variables x, y, u . Set for $b \in X$

$$\begin{aligned} r_{\Sigma_1}^\sigma(i, b) &\simeq \text{the } <_L\text{-least } c \in L_\sigma \text{ such that } L_\sigma \models \theta_i((c)_0, (c)_1; b) \\ h_{\Sigma_1}^\sigma(i, b) &\simeq (r_{\Sigma_1}^\sigma(i, b))_0 \\ \text{Hull}_{\Sigma_1}^\sigma(X) &= \text{rng}(h_{\Sigma_1}^\sigma) = \{h_{\Sigma_1}^\sigma(i, b) \in L_\sigma : i \in \omega, b \in X\} \end{aligned}$$

Then $L_\sigma \models \exists x \exists y \theta_i(x, y; b) \rightarrow h_{\Sigma_1}^\sigma(i, b) \downarrow \& \exists y \theta_i(h_{\Sigma_1}^\sigma(i, b), y; b)$.

Proposition 2.2 Assume that X is a set in L_σ . Then $r_{\Sigma_1}^\sigma$ and $h_{\Sigma_1}^\sigma$ are partial $\Delta_1(L_\sigma)$ -maps such that the domain of $h_{\Sigma_1}^\sigma$ is a $\Sigma_1(L_\sigma)$ -subset of $\omega \times X$. Therefore its range $\text{Hull}_{\Sigma_1}^\sigma(X)$ is a $\Sigma_1(L_\sigma)$ -subset of L_σ .

Proposition 2.3 Let $Y = \text{Hull}_{\Sigma_1}^\sigma(X)$. For any $\Sigma_1(Y)$ -sentence $\varphi(\bar{a})$ with parameters \bar{a} from Y $L_\sigma \models \varphi(\bar{a}) \Leftrightarrow Y \models \varphi(\bar{a})$. Namely $Y \prec_{\Sigma_1} L_\sigma$.

Definition 2.4 (Mostowski collapsing function F)

By Proposition 2.3 and the Condensation Lemma we have an isomorphism (Mostowski collapsing function)

$$F : \text{Hull}_{\Sigma_1}^\sigma(X) \leftrightarrow L_\gamma$$

for an ordinal $\gamma \leq \sigma$ such that $F \upharpoonright Y = id \upharpoonright Y$ for any transitive $Y \subset \text{Hull}_{\Sigma_1}^\sigma(X)$.

Let us denote, though $\sigma \notin \text{dom}(F) = \text{Hull}_{\Sigma_1}^\sigma(X)$

$$F(\sigma) := \gamma.$$

Also for the above Mostowski collapsing map F let

$$F^{\Sigma_1}(x; \sigma, X) := F(x).$$

The inverse $G := F^{-1}$ of F is a Σ_1 -elementary embedding from $L_{F(\sigma)}$ to L_σ .

Proposition 2.5 Let $L_\sigma \models \text{KP}\omega + \Sigma_1\text{-Collection}$. Then for $\kappa \leq \sigma$, $\{(x, y) : x < \kappa \& y = \min\{y < \kappa : \text{Hull}_{\Sigma_1}^\sigma(x \cup \{\kappa\}) \cap \kappa \subset y\}\}$ is a $\text{Bool}(\Sigma_1(L_\sigma))$ -predicate on κ , and hence a set in L_σ if $\kappa < \sigma$ and $L_\sigma \models \Sigma_1\text{-Separation}$.

$F_{x \cup \{\kappa\}}^{\Sigma_1}(y)$ denotes the Mostowski collapse $F^{\Sigma_1}(y; \sigma, x \cup \{\kappa\})$.

Theorem 2.6 Let σ be an ordinal such that $L_\sigma \models \text{KP}\omega + \Sigma_1\text{-Separation}$, and $\omega \leq \alpha < \kappa < \sigma$ with α a multiplicative principal number and κ a limit ordinal. Then the following conditions are mutually equivalent:

1. $L_\sigma \models {}^\alpha\kappa \subset L_\kappa$.
2. $L_\sigma \models \alpha < cf(\kappa)$.
3. There exists an ordinal x such that $\alpha < x < \kappa$, $\text{Hull}_{\Sigma_1}^\sigma(x \cup \{\kappa\}) \cap \kappa \subset x$ and $F_{x \cup \{\kappa\}}^{\Sigma_1}(\sigma) < \kappa$.
4. For the Mostowski collapse $F_{x \cup \{\kappa\}}^{\Sigma_1}(y)$, there exists an ordinal x such that $\alpha < x = F_{x \cup \{\kappa\}}^{\Sigma_1}(\kappa) < F_{x \cup \{\kappa\}}^{\Sigma_1}(\sigma) < \kappa$, and for any Σ_1 -formula φ and any $a \in L_x$, $L_\sigma \models \varphi[\kappa, a] \rightarrow L_{F_{x \cup \{\kappa\}}^{\Sigma_1}(\sigma)} \models \varphi[x, a]$ holds.

Definition 2.7 $T_1 := \text{KP}\omega + (V = L) + (\Pi_1\text{-Collection}) + (\omega_1)$ denotes an extension of the Kripke-Platek set theory with the axioms of infinity, constructibility, $\Pi_1\text{-Collection}$ and the following axiom:

$$(\omega_1) \exists \kappa \forall \alpha < \kappa \exists \beta, \gamma < \kappa [\alpha < \beta < \gamma \wedge L_\gamma = \text{rng}(F_{\beta \cup \{\kappa\}}^{\Sigma_1}) \wedge \text{Hull}_{\Sigma_1}(\beta \cup \{\kappa\}) \cap \kappa \subset \beta]$$

where $F_{\beta \cup \{\kappa\}}^{\Sigma_1} : \text{Hull}_{\Sigma_1}(\beta \cup \{\kappa\}) \rightarrow L_\gamma$ is the Mostowski collapsing map, and $\text{Hull}_{\Sigma_1}(x)$ is the Σ_1 -Skolem hull of sets x in the universe.

From Theorem 2.6 we see that $T_1 \vdash \exists \kappa \forall \alpha < \kappa (\alpha < cf(\kappa))$.

3 A theory equivalent to T_1

Referring Theorem 2.6 let us interpret T_1 to another theory. The base language here is $\{\in\}$.

Let ρ_0 denotes the least ordinal above the least uncountable ordinal ω_1 such that $L_{\rho_0} \models (\Pi_1\text{-Collection})$. $F_X(x) := F^{\Sigma_1}(x; \rho_0, X)$ and $\text{Hull}(X) := \text{Hull}_{\Sigma_1}^{\rho_0}(X)$.

The predicate P is intended to denote the relation $P(x, y)$ iff $x = F_{x \cup \{\omega_1\}}^{\Sigma_1}(\omega_1)$ and $y = F_{x \cup \{\omega_1\}}^{\Sigma_1}(\rho_0)$. Also the predicate $P_{\rho_0}(x)$ is intended to denote the relation $P_{\rho_0}(x)$ iff $x = F_x(\rho_0)$.

Definition 3.1 $T(\omega_1)$ denotes the set theory defined as follows.

1. Its language is $\{\in, P, P_{\rho_0}, \omega_1\}$ for a binary predicate P , a unary predicate P_{ρ_0} and an individual constant ω_1 .

2. Its axioms are obtained from those of KP ω + (Π_1 -Collection) in the expanded language¹, the axiom of constructibility $V = L$ together with the axiom schema saying that ω_1 is an uncountable regular ordinal, cf. (2) and (1), and if $P(x, y)$ then x is a critical point of the Σ_1 -elementary embedding from $L_y \cong \text{Hull}(x \cup \{\omega_1\})$ to the universe L_{ρ_0} , cf. (1), and if $P_{\rho_0}(x)$ then x is a critical point of the Σ_1 -elementary embedding from $L_x \cong \text{Hull}(x)$ to the universe L_{ρ_0} , cf.(3): for a formula φ and an ordinal α , φ^α denotes the result of restricting every unbounded quantifier $\exists z, \forall z$ in φ to $\exists z \in L_\alpha, \forall z \in L_\alpha$.

(a) $x \in Ord$ is a Δ_0 -formula saying that ‘ x is an ordinal’.

$$(\omega < \omega_1 \in Ord), (P(x, y) \rightarrow \{x, y\} \subset Ord \wedge x < y < \omega_1) \text{ and } (P_{\rho_0}(x) \rightarrow x \in Ord).$$

(b)

$$P(x, y) \rightarrow a \in L_x \rightarrow \varphi[\omega_1, a] \rightarrow \varphi^y[x, a] \quad (1)$$

for any Σ_1 -formula φ in the language $\{\in\}$.

(c)

$$a \in Ord \cap \omega_1 \rightarrow \exists x, y \in Ord \cap \omega_1 [a < x \wedge P(x, y)] \quad (2)$$

(d)

$$P_{\rho_0}(x) \rightarrow a \in L_x \rightarrow \varphi[a] \rightarrow \varphi^x[a] \quad (3)$$

for any Σ_1 -formula φ in the language $\{\in\}$.

(e)

$$a \in Ord \rightarrow \exists x \in Ord [a < x \wedge P_{\rho_0}(x)] \quad (4)$$

Remark. Though the axioms (3) and (4) for the Π_1 -definable predicate $P_{\rho_0}(x)$ are derivable from Π_1 -Collection, the primitive predicate symbol $P_{\rho_0}(x)$ is useful for our prof-theoretic study, cf. the proof of Lemma 5.20 below.

Lemma 3.2 $T(\omega_1)$ is a conservative extension of the set theory T_1 .

Proof. First consider the axioms of T_1 in $T(\omega_1)$. The axiom (ω_1) follows from (1). Hence we have shown that T_1 is contained in $T(\omega_1)$.

Next we show that $T(\omega_1)$ is interpretable in T_1 . Let κ be an ordinal in the axiom (ω_1) . Interpret the predicate $P(x, y) \leftrightarrow \{x, y\} \subset Ord \wedge (\text{Hull}(x \cup \{\kappa\}) \cap \kappa \subset x) \wedge (y = \sup\{F_{x \cup \{\kappa\}}(a) : a \in \text{Hull}(x \cup \{\kappa\})\})$. We see from Theorem 2.6 that the interpreted (1) and (2) are provable in T_1 .

It remains to show the interpreted (3) and (4) in T_1 . It suffices to show that given an ordinal α , there exists an ordinal $x > \alpha$ such that $\text{Hull}(x) \cap Ord \subset x$.

First we show that for any α there exists a β such that $\text{Hull}(\alpha) \cap Ord \subset \beta$. By Proposition 2.2 let $h_{\Sigma_1}^{\rho_0}$ be the Δ_1 -surjection from the Σ_1 -subset $\text{dom}(h_{\Sigma_1}^{\rho_0})$ of

¹ This means that the predicates P, P_{ρ_0} do not occur in Δ_0 -formulae for Δ_0 -Separation and Π_1 -formulae Π_1 -Collection.

$\omega \times \alpha$ to $\text{Hull}(\alpha)$, which is a Σ_1 -class. From Σ_1 -Separation we see that $\text{dom}(h_{\Sigma_1}^{\rho_0})$ is a set. Hence by Σ_1 -Collection, $\text{Hull}(\alpha) = \text{rng}(h_{\Sigma_1}^{\rho_0})$ is a set. Therefore the ordinal $\sup(\text{Hull}(\alpha) \cap \text{Ord})$ exists in the universe.

As in Proposition 2.5 we see that $X = \{(\alpha, \beta) : \beta = \min\{\beta \in \text{Ord} : \text{Hull}(\alpha) \cap \text{Ord} \subset \beta\}\}$ is a set in L_{ρ_0} as follows. Let $\varphi(\beta)$ be the Π_1 -predicate $\varphi(\beta) :\Leftrightarrow \forall \gamma \in \text{Ord}[\gamma \in \text{Hull}(\alpha) \rightarrow \gamma \in \beta]$. Then $\beta = \min\{\beta : \text{Hull}(\alpha) \cap \text{Ord} \subset \beta\}$ iff $\varphi(\beta) \wedge \forall \gamma < \beta \neg \varphi(\gamma)$, which is $\text{Bool}(\Sigma_1(L_{\rho_0}))$ by Π_0 -Collection. Hence X is a set in L_{ρ_0} .

Define recursively ordinals $\{x_n\}_n$ as follows. $x_0 = \alpha + 1$, and x_{n+1} is defined to be the least ordinal x_{n+1} such that $\text{Hull}(x_n) \cap \text{Ord} \subset x_{n+1}$, i.e., $(x_n, x_{n+1}) \in X$. We see inductively that such an ordinal x_n exists. Moreover $n \mapsto x_n$ is a Δ_1 -map. Then $x = \sup_n x_n < \rho_0$ is a desired one. \square

4 Ordinals for ω_1

For our proof-theoretic analysis of T_1 , we need to talk about ‘ordinals’ less than the next epsilon number to the order type of the class of ordinals inside T_1 . Let $\text{Ord}^\varepsilon \subset V$ and $<^\varepsilon$ be Δ -predicates such that for any transitive and wellfounded model V of $\text{KP}\omega$, $<^\varepsilon$ is a well ordering of type ε_{ρ_0+1} on Ord^ε for the order type ρ_0 of the class Ord in V . $<^\varepsilon$ is seen to be a canonical ordering as stated in the following Proposition 4.1.

Proposition 4.1 1. $\text{KP}\omega$ proves the fact that $<^\varepsilon$ is a linear ordering.

2. For any formula φ and each $n < \omega$,

$$\text{KP}\omega \vdash \forall x \in \text{Ord}^\varepsilon (\forall y <^\varepsilon x \varphi(y) \rightarrow \varphi(x)) \rightarrow \forall x <^\varepsilon \omega_n(\rho_0 + 1)\varphi(x) \quad (5)$$

In what follows of this section we work in T_1 . For simplicity let us identify the code $x \in \text{Ord}^\varepsilon$ with the ‘ordinal’ coded by x , and $<^\varepsilon$ is denoted by $<$ when no confusion likely occurs. Note that the ordinal ρ_0 is the order type of the class of ordinals in the intended model L_{ρ_0} of T_1 . Define simultaneously the classes $\mathcal{H}_\alpha(X) \subset L_{\rho_0} \cup \varepsilon_{\rho_0+1}$ and the ordinals $\Psi_{\omega_1}(\alpha)$ and $\Psi_{\rho_0}(\alpha)$ for $\alpha <^\varepsilon \varepsilon_{\rho_0+1}$ and sets $X \subset L_{\omega_1}$ as follows. We see that $\mathcal{H}_\alpha(X)$ and $\Psi_\kappa(\alpha)$ ($\kappa \in \{\omega_1, \rho_0\}$) are (first-order) definable as a fixed point in T_1 , cf. Proposition 4.4.

Recall that $\text{Hull}(X) = \text{Hull}_{\Sigma_1}^{\rho_0}(X) \subset L_{\rho_0}$ and $F_X(x) = F^{\Sigma_1}(x; \rho_0, X)$ with $F_X : \text{Hull}(X) \rightarrow L_\gamma$ for $X \subset L_{\rho_0}$ and a $F_X(\rho_0) = \gamma \leq \rho_0$.

Definition 4.2 $\mathcal{H}_\alpha(X)$ is the Skolem hull of $\{0, \omega_1, \rho_0\} \cup X$ under the functions $+, \alpha \mapsto \omega^\alpha, \Psi_{\omega_1} \upharpoonright \alpha, \Psi_{\rho_0} \upharpoonright \alpha$, the Σ_1 -definability, and the Mostowski collapsing functions $(x, d) \mapsto F_{x \cup \{\omega_1\}}(d)$ ($\text{Hull}(x \cup \{\omega_1\}) \cap \omega_1 \subset x$) and $d \mapsto F_x(d)$ ($\text{Hull}(x) \cap \rho_0 \subset x$).

1. $\{0, \omega_1, \rho_0\} \cup X \subset \mathcal{H}_\alpha(X)$.
2. $x, y \in \mathcal{H}_\alpha(X) \Rightarrow x + y, \omega^x \in \mathcal{H}_\alpha(X)$.
3. $\gamma \in \mathcal{H}_\alpha(X) \cap \alpha \Rightarrow \Psi_\kappa(\gamma) \in \mathcal{H}_\alpha(X)$ for $\kappa \in \{\omega_1, \rho_0\}$.

4. $\text{Hull}(\mathcal{H}_\alpha(X) \cap L_{\rho_0}) \subset \mathcal{H}_\alpha(X)$.

Namely for any Σ_1 -formula $\varphi[x, \vec{y}]$ in the language $\{\in\}$ and parameters $\vec{a} \subset \mathcal{H}_\alpha(X) \cap L_{\rho_0}$, if $b \in L_{\rho_0}$, $(L_{\rho_0}, \in) \models \varphi[b, \vec{a}]$ and $(L_{\rho_0}, \in) \models \exists!x \varphi[x, \vec{a}]$, then $b \in \mathcal{H}_\alpha(X)$.

5. If $x \in \mathcal{H}_\alpha(X) \cap \omega_1$ with $\text{Hull}(x \cup \{\omega_1\}) \cap \omega_1 \subset x$, and $d \in (\text{Hull}(x \cup \{\omega_1\}) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$, then $F_{x \cup \{\omega_1\}}(d) \in \mathcal{H}_\alpha(X)$.
6. If $x \in \mathcal{H}_\alpha(X) \cap \rho_0$ with $\text{Hull}(x) \cap \rho_0 \subset x$, and $d \in (\text{Hull}(x) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$, then $F_x(d) \in \mathcal{H}_\alpha(X)$.

For $\kappa \in \{\omega_1, \rho_0\}$

$$\Psi_\kappa(\alpha) := \min\{\beta \leq \kappa : \mathcal{H}_\alpha(\beta) \cap \kappa \subset \beta\}.$$

The ordinal $\Psi_\kappa(\alpha)$ is well defined and $\Psi_\kappa(\alpha) \leq \kappa$ for $\kappa \in \{\omega_1, \rho_0\}$.

- Proposition 4.3**
1. $\mathcal{H}_\alpha(X)$ is closed under Σ_1 -definability: $\vec{a} \subset \mathcal{H}_\alpha(X) \cap L_{\rho_0} \Rightarrow \text{Hull}(\vec{a}) \subset \mathcal{H}_\alpha(X)$.
 2. $\text{Hull}(\Psi_{\omega_1}(\alpha) \cup \{\omega_1\}) \cap \omega_1 = \Psi_{\omega_1}(\alpha)$ and $\text{Hull}(\Psi_{\rho_0}(\alpha)) \cap \rho_0 = \Psi_{\rho_0}(\alpha) > \omega_1$
 3. $\mathcal{H}_\alpha(X)$ is closed under the Veblen function φ on ρ_0 , $x, y \in \mathcal{H}_\alpha(X) \cap \rho_0 \Rightarrow \varphi xy \in \mathcal{H}_\alpha(X)$.
 4. If $x \in \mathcal{H}_\alpha(X) \cap \omega_1$, $\text{Hull}(x \cup \{\omega_1\}) \cap \omega_1 \subset x$, and $\delta \in (\text{Hull}(x \cup \{\omega_1\}) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$, then $F_{x \cup \{\omega_1\}}(\delta) \in \mathcal{H}_\alpha(X)$.
 5. If $x \in \mathcal{H}_\alpha(X) \cap \rho_0$, $\text{Hull}(x) \cap \rho_0 \subset x$, and $\delta \in (\text{Hull}(x) \cap \mathcal{H}_\alpha(X)) \cup \{\rho_0\}$, then $F_x(\delta) \in \mathcal{H}_\alpha(X)$.

The following Proposition 4.4 is easy to see.

Proposition 4.4 Both of $x = \mathcal{H}_\alpha(X)$ and $y = \Psi_\kappa(\alpha)$ ($\kappa \in \{\omega_1, \rho_0\}$) are Σ_2 -predicates as fixed points in $\text{KP}\omega$.

Lemma 4.5 For each $n < \omega$,

$$T_1 \vdash \forall \alpha < \omega_{n+1}(\rho_0 + 1) \forall \kappa \in \{\omega_1, \rho_0\} \exists x < \kappa [x = \Psi_\kappa(\alpha)].$$

Proof. Let $\kappa \in \{\omega_1, \rho_0\}$. By Proposition 4.4 both $x = \mathcal{H}_\alpha(\beta)$ and $y = \Psi_\kappa(\alpha)$ are Σ_2 -predicates. We show that $A(\alpha) :\Leftrightarrow \forall \beta < \rho_0 \exists x [x = \mathcal{H}_\alpha(\beta)] \wedge \forall \kappa \in \{\omega_1, \rho_0\} \exists \beta < \kappa [\Psi_\kappa(\alpha) = \beta]$ is progressive. Then $\forall \alpha < \omega_{n+1}(\rho_0 + 1) \forall \kappa \in \{\omega_1, \rho_0\} \exists x < \kappa [x = \Psi_\kappa(\alpha)]$ will follow from transfinite induction up to $\omega_{n+1}(\rho_0 + 1)$, cf. (5) in Proposition 4.1.

Assume $\forall \gamma < \alpha A(\gamma)$ as our IH. Since $\text{dom}(h_{\Sigma_1}^{\rho_0})$ is a Σ_1 -subset of $\omega \times \beta$ for $\beta < \rho_0$, it is a set by Σ_1 -Separation. Then so is the image $\text{Hull}(\beta)$ of the Δ_1 -map $h_{\Sigma_1}^{\rho_0}$. Hence $\forall \beta < \rho_0 \exists h [h = \text{Hull}(\beta)]$.

We see from this, IH and Σ_2 -Collection that $\forall\beta < \rho_0 \exists x[x = \mathcal{H}_\alpha(\beta)] = \bigcup_m \mathcal{H}_\alpha^m(\beta)$, where $\mathcal{H}_\alpha^m(\beta)$ is an m -th stage of the construction of $\mathcal{H}_\alpha(\beta)$ such that $x = \mathcal{H}_\alpha^m(\beta)$ is a Σ_2 -predicate.

Define recursively ordinals $\{\beta_m\}_m$ for $\kappa \in \{\omega_1, \rho_0\}$ as follows. $\beta_0 = 0$, and β_{m+1} is defined to be the least ordinal $\beta_{m+1} \leq \kappa$ such that $\mathcal{H}_\alpha(\beta_m) \cap \kappa \subset \beta_{m+1}$.

We see inductively that $\beta_m < \kappa$ using the regularity of κ and the facts that $\forall\beta < \kappa \exists x[x = \mathcal{H}_\alpha(\beta) \wedge \text{card}(x) < \kappa]$, where $\text{card}(x) < \kappa$ designates that there exists a surjection $f : \gamma \rightarrow x$ for a $\gamma < \kappa$ and $f \in L_{\rho_0}$. Moreover $m \mapsto \beta_m$ is a Σ_2 -map. Therefore $\beta = \sup_m \beta_m < \kappa$ enjoys $\mathcal{H}_\alpha(\beta) \cap \kappa \subset \beta$. \square

5 Operator controlled derivations for T_1

5.1 An intuitionistic fixed point theory $\text{FiX}^i(T_1)$

Let us introduce an intuitionistic fixed point theory $\text{FiX}^i(T_1)$ over the set theory T_1 . Fix an X -strictly positive formula $\mathcal{Q}(X, x)$ in the language $\{\in, =, X\}$ with an extra unary predicate symbol X . In $\mathcal{Q}(X, x)$ the predicate symbol X occurs only strictly positive. The language of $\text{FiX}^i(T_1)$ is $\{\in, =, Q\}$ with a fresh unary predicate symbol Q . The axioms in $\text{FiX}^i(T_1)$ consist of the following:

1. All derivable sentences in T_1 in the language $\{\in\}$.
2. Induction schema for any formula φ in $\{\in, =, Q\}$:

$$\forall x(\forall y \in x \varphi(y) \rightarrow \varphi(x)) \rightarrow \forall x \varphi(x).$$
3. Fixed point axiom: $\forall x[Q(x) \leftrightarrow \mathcal{Q}(Q, x)]$.

The underlying logic in $\text{FiX}^i(T_1)$ is defined to be the intuitionistic first-order predicate logic with equality.

Lemma 5.1 *Let $<^\varepsilon$ denote a Δ_1 -predicate defined in section 4. For each $n < \omega$ and each formula φ in $\{\in, =, Q\}$,*

$$\text{FiX}^i(T_1) \vdash \forall x(\forall y <^\varepsilon x \varphi(y) \rightarrow \varphi(x)) \rightarrow \forall x <^\varepsilon \omega_n(\rho_0 + 1)\varphi(x).$$

The following Theorem 5.2 is shown in [2].

Theorem 5.2 *$\text{FiX}^i(T_1)$ is a conservative extension of T_1 .*

5.2 Classes of formulae

In this subsection we work in T_1 .

The language \mathcal{L}_c is obtained from $\{\in, P, P_{\rho_0}, \omega_1\}$ by adding names (individual constants) c_a of each set $a \in L_{\rho_0}$. c_a is identified with a . A *term* in \mathcal{L}_c is either a variable or a constant in L_{ρ_0} . Formulae in this language are defined in the next definition. Formulae are assumed to be in negation normal form.

Definition 5.3 1. Let t_1, \dots, t_m be terms. For each m -ary predicate constant $R \in \{\in, P, P_{\rho_0}\}$ $R(t_1, \dots, t_m)$ and $\neg R(t_1, \dots, t_m)$ are formulae, where $m = 1, 2$. These are called *literals*.

2. If A and B are formulae, then so are $A \wedge B$ and $A \vee B$.
3. Let t be a term. If A is a formula and the variable x does not occur in t , then $\exists x \in t A$ and $\forall x \in t A$ are formulae. $\exists x \in t A$, $\forall x \in t$ are *bounded quantifiers*.
4. If A is a formula and x a variable, then $\exists x A$ and $\forall x A$ are formulae. Unbounded quantifiers $\exists x, \forall x$ are denoted by $\exists x \in L_{\rho_0}, \forall x \in L_{\rho_0}$, resp.

For formulae A in \mathcal{L}_c , $\text{qk}(A)$ denotes the finite set of sets $a \in L_{\rho_0}$ which are bounds of bounded quantifiers $\exists x \in a, \forall x \in a$ occurring in A . Moreover $\text{k}(A)$ denotes the set of sets occurring in A . $\text{k}(A)$ is defined to include bounds of bounded quantifiers. By definition we set $0 \in \text{qk}(A)$. Thus $0 \in \text{qk}(A) \subset \text{k}(A) \subset L_{\rho_0}$.

Definition 5.4 1. $\text{k}(\neg A) = \text{k}(A)$ and similarly for qk .

2. $\text{qk}(M) = \{0\}$ for any literal M .
3. $\text{k}(Q(t_1, \dots, t_m)) = (\{t_1, \dots, t_m\} \cap L_{\rho_0}) \cup \{0\}$ for literals $Q(t_1, \dots, t_m)$ with predicates Q in the set $\{\in, P, P_{\rho_0}\}$.
4. $\text{k}(A_0 \vee A_1) = \text{k}(A_0) \cup \text{k}(A_1)$ and similarly for qk .
5. For unbounded quantifiers, $\text{k}(\exists x A(x)) = \text{k}(A(x))$ and similarly for qk .
6. For bounded quantifiers with $a \in L_{\rho_0}$, $\text{k}(\exists x \in a A(x)) = \{a\} \cup \text{k}(A(x))$ and similarly for qk .
7. For variables y , $\text{k}(\exists x \in y A(x)) = \text{k}(A(x))$ and similarly for qk .
8. For sets Γ of formulae $\text{k}(\Gamma) := \bigcup \{\text{k}(A) : A \in \Gamma\}$.

For example $\text{qk}(\exists x \in a A(x)) = \{a\} \cup \text{qk}(A(x))$ if $a \in L_{\rho_0}$.

Definition 5.5 For $a \in L_{\rho_0} \cup \{L_{\rho_0}\}$, $\text{rk}_L(a)$ denotes the *L-rank* of a .

$$\text{rk}_L(a) := \begin{cases} \min\{\alpha \in \text{Ord} : a \in L_{\alpha+1}\} & a \in L_{\rho_0} \\ \rho_0 & a = L_{\rho_0} \end{cases}$$

Definition 5.6 1. $A \in \Delta_0$ iff there exists a Δ_0 -formula $\theta[\vec{x}]$ in the language $\{\in\}$ and terms \vec{t} in \mathcal{L}_c such that $A \equiv \theta[\vec{t}]$. This means that A is bounded, and the predicates P, P_{ρ_0} do not occur in A .

2. Putting $\Sigma_0 := \Pi_0 := \Delta_0$, the classes Σ_m and Π_m of formulae in the language \mathcal{L}_c are defined as usual, where by definition $\Sigma_m \cup \Pi_m \subset \Sigma_{m+1} \cap \Pi_{m+1}$.

Each formula in $\Sigma_m \cup \Pi_m$ is in prenex normal form with alternating unbounded quantifiers and Δ_0 -matrix.

3. The set $\Sigma^{\Sigma_n}(\lambda)$ of sentences is defined recursively as follows. Let $\{a, b, c\} \subset L_{\rho_0}$ and $d \in L_{\rho_0} \cup \{L_{\rho_0}\}$.
 - (a) Each Σ_n -sentence is in $\Sigma^{\Sigma_n}(\lambda)$.
 - (b) Each literal including $Reg(a), P(a, b, c), P_{I,n}(a)$ and its negation is in $\Sigma^{\Sigma_n}(\lambda)$.
 - (c) $\Sigma^{\Sigma_n}(\lambda)$ is closed under propositional connectives \vee, \wedge .
 - (d) Suppose $\forall x \in d A(x) \notin \Delta_0$. Then $\forall x \in d A(x) \in \Sigma^{\Sigma_n}(\lambda)$ iff $A(\emptyset) \in \Sigma^{\Sigma_n}(\lambda)$ and $\text{rk}_L(d) < \lambda$.
 - (e) Suppose $\exists x \in d A(x) \notin \Delta_0$. Then $\exists x \in d A(x) \in \Sigma^{\Sigma_n}(\lambda)$ iff $A(\emptyset) \in \Sigma^{\Sigma_n}(\lambda)$ and $\text{rk}_L(d) \leq \lambda$.
4. For a Σ_1 -formula $A(x)$, $\exists x \in P_{\rho_0} A(x)$ is a $\Sigma_1(P_{\rho_0})$ -formula.

Note that the predicates P, P_{ρ_0} do not occur in Σ_m -formulae.

Definition 5.7 Let us extend the domain $\text{dom}(F_x) = \text{Hull}(x)$ of the Mostowski collapse to formulae.

$$\text{dom}(F_x) = \{A \in \Sigma_1 \cup \Pi_1 : \mathbf{k}(A) \subset \text{Hull}(x)\}.$$

For $A \in \text{dom}(F_x)$, $F_x''A$ denotes the result of replacing each constant $c \in L_{\rho_0}$ by $F_x(c)$, each unbounded existential quantifier $\exists z \in L_{\rho_0}$ by $\exists z \in L_{F_x(\rho_0)}$, and each unbounded universal quantifier $\forall z \in L_{\rho_0}$ by $\forall z \in L_{F_x(\rho_0)}$.

For sequent, i.e., finite set of sentences $\Gamma \subset \text{dom}(F_x)$, put $F_x''\Gamma = \{F_x''A : A \in \Gamma\}$.

The assignment of disjunctions $A \simeq \bigvee(A_\iota)_{\iota \in J}$ or conjunctions $A \simeq \bigwedge(A_\iota)_{\iota \in J}$ to sentences A is defined as in [3] except for $\Sigma_1 \cup \Pi_1$ -sentences.

Definition 5.8 1. If M is one of the literals $a \in b, a \notin b$, then for $J := 0$

$$M : \simeq \begin{cases} \bigvee(A_\iota)_{\iota \in J} & \text{if } M \text{ is false (in } L) \\ \bigwedge(A_\iota)_{\iota \in J} & \text{if } M \text{ is true} \end{cases}$$

2. $(A_0 \vee A_1) : \simeq \bigvee(A_\iota)_{\iota \in J}$ and $(A_0 \wedge A_1) : \simeq \bigwedge(A_\iota)_{\iota \in J}$ for $J := 2$.

3. $P(b, c) : \simeq \bigvee(0 \notin 0)_{\iota \in J}$ and $\neg P(b, c) : \simeq \bigwedge(0 \in 0)_{\iota \in J}$ with

$$J := \begin{cases} 1 & \text{if } \exists \alpha[b = \Psi_{\omega_1}(\alpha) \& c = F_{b \cup \{\omega_1\}}(\rho_0)] \\ 0 & \text{otherwise} \end{cases}.$$

4. $P_{\rho_0}(a) : \simeq \bigvee(0 \notin 0)_{\iota \in J}$ and $\neg P_{\rho_0}(a) : \simeq \bigwedge(0 \in 0)_{\iota \in J}$ with

$$J := \begin{cases} 1 & \text{if } \exists \alpha[a = \Psi_{\rho_0}(\alpha)] \\ 0 & \text{otherwise} \end{cases}.$$

5. Let $\exists z \in b \theta[z] \in \Sigma_0$ for $b \in L_{\rho_0} \cup \{L_{\rho_0}\}$. Then for the set

$$d := \mu z \in b \theta[z] := \begin{cases} \min_{<_L} \{d : d \in b \wedge \theta[d]\} & \text{if } \exists z \in b \theta[z] \\ 0 & \text{otherwise} \end{cases}$$

with a canonical well ordering $<_L$ on L , and $J = \{d\}$

$$\begin{aligned} \exists z \in b \theta[z] &:= \bigvee (d \in b \wedge \theta[d])_{d \in J} \\ \forall z \in b \neg \theta[z] &:= \bigwedge (d \in b \rightarrow \neg \theta[d])_{d \in J} \end{aligned}$$

where $d \in b$ denotes a true literal, e.g., $d \notin d$ when $b = L_{\rho_0}$.

6. For a $\Sigma_1(P_{\rho_0})$ -sentence $\exists x \in P_{\rho_0} A(x)$,

$$\begin{aligned} \exists x \in P_{\rho_0} A(x) &\simeq \bigvee (A(a))_{a \in J} \\ \forall x \in P_{\rho_0} \neg A(x) &\simeq \bigwedge (\neg A(a))_{a \in J} \\ \text{with } J &= \{a : \exists \alpha (a = \Psi_{\rho_0}(\alpha))\} \end{aligned}$$

7. Otherwise set for $a \in L_{\rho_0} \cup \{L_{\rho_0}\}$ and $J := \{b : b \in a\}$

$$\exists x \in a A(x) := \bigvee (A(b))_{b \in J} \text{ and } \forall x \in a A(x) := \bigwedge (A(b))_{b \in J}.$$

The rank $\text{rk}(A)$ of sentences A is defined by recursion on the number of symbols occurring in A .

Definition 5.9 1. $\text{rk}(\neg A) := \text{rk}(A)$.

2. $\text{rk}(a \in b) := 0$.

3. $\text{rk}(P(b, c)) := \text{rk}(P_{\rho_0}(a)) := 1$.

4. $\text{rk}(A_0 \vee A_1) := \max\{\text{rk}(A_0), \text{rk}(A_1)\} + 1$.

5. $\text{rk}(\exists x \in a A(x)) := \max\{\omega \alpha, \text{rk}(A(\emptyset)) + 1\}$ for $\alpha = \text{rk}_L(a)$.

6. $\text{rk}(\exists x \in P_{\rho_0} A(x)) = \rho_0$.

Proposition 5.10 Let $A \simeq \bigvee (A_\iota)_{\iota \in J}$ or $A \simeq \bigwedge (A_\iota)_{\iota \in J}$.

1. $\forall \iota \in J (\mathbf{k}(A_\iota) \subset \mathbf{k}(A) \cup \{\iota\})$.

2. $A \in \Sigma^{\Sigma_n}(\lambda) \Rightarrow \forall \iota \in J (A_\iota \in \Sigma^{\Sigma_n}(\lambda))$.

3. For an ordinal $\lambda \leq \rho_0$ with $\omega \lambda = \lambda$, $\text{rk}(A) < \lambda \Rightarrow A \in \Sigma^{\Sigma_n}(\lambda)$.

4. $\text{rk}(A) < \rho_0 + \omega$.

5. $\text{rk}(A) \in \{\omega \text{rk}_L(a) + i : a \in \mathbf{qk}(A) \cup \{\rho_0\}, i \in \omega\} \subset \text{Hull}(\mathbf{k}(A))$.

6. $\forall \iota \in J (\text{rk}(A_\iota) < \text{rk}(A))$.

5.3 Operator controlled derivations

In the remaining parts of this section we work in the intuitionistic fixed point theory $\text{FiX}^i(T_1)$.

Sequents are finite sets of sentences, and inference rules are formulated in one-sided sequent calculus. In what follows by an operator we mean an $\mathcal{H}_\gamma[\Theta]$ for a finite set Θ of sets.

Definition 5.11 Define a relation $(\mathcal{H}, \kappa) \vdash_b^a \Gamma$ as follows.

$(\mathcal{H}, \kappa) \vdash_b^a \Gamma$ holds if

$$\{a\} \cup \kappa(\Gamma) \subset \mathcal{H} := \mathcal{H}(\emptyset) \quad (6)$$

and one of the following cases holds:

(\vee) $A \simeq \bigvee \{A_\iota : \iota \in J\}$, $A \in \Gamma$ and there exist $\iota \in J$ and $a(\iota) < a$ such that

$$\text{rk}_L(\iota) < \kappa \Rightarrow \text{rk}_L(\iota) < a \quad (7)$$

and $(\mathcal{H}, \kappa) \vdash_b^{a(\iota)} \Gamma, A_\iota$.

(\wedge) $A \simeq \bigwedge \{A_\iota : \iota \in J\}$, $A \in \Gamma$ and for every $\iota \in J$ there exists an $a(\iota) < a$ such that $(\mathcal{H}[\{\iota\}], \kappa) \vdash_b^{a(\iota)} \Gamma, A_\iota$.

(cut) There exist $a_0 < a$ and C such that $\text{rk}(C) < b$ and $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Gamma, \neg C$ and $(\mathcal{H}, \kappa) \vdash_b^{a_0} C, \Gamma$.

(P) There exists $\alpha < \omega_1$ such that $(\exists x, y < \omega_1 [\alpha < x \wedge P(x, y)]) \in \Gamma$.

(F_x) $x = \Psi_{\omega_1}(\beta) \in \mathcal{H}$ for a β and there exist $a_0 < a$, $\Gamma_0 \subset \Sigma_1$ and Λ such that $\kappa(\Gamma_0) \subset \text{Hull}(x \cup \{\omega_1\})$, $\Gamma = \Lambda \cup (F_{x \cup \{\omega_1\}}'' \Gamma_0)$ and $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Lambda, \Gamma_0$, where $F_{x \cup \{\omega_1\}}$ denotes the Mostowski collapse $F_{x \cup \{\omega_1\}} : \text{Hull}(x \cup \{\omega_1\}) \leftrightarrow L_{F_{x \cup \{\omega_1\}}(\rho_0)}$.

(P_{ρ₀}) There exists $\alpha < \rho_0$ such that $(\exists x < \rho_0 [\alpha < x \wedge P_{\rho_0}(x)]) \in \Gamma$.

(F_x) $x = \Psi_{\rho_0}(\beta) \in \mathcal{H}$ for a β and there exist $a_0 < a$, $\Gamma_0 \subset \Sigma_1$ and Λ such that $\kappa(\Gamma_0) \subset \text{Hull}(x)$, $\Gamma = \Lambda \cup (F_x'' \Gamma_0)$ and $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Lambda, \Gamma_0$, where F_x denotes the Mostowski collapse $F_x : \text{Hull}(x) \leftrightarrow L_{F_x(\rho_0)}$.

(Ref) $b \geq \rho_0$, and there exist an ordinal $a_0 < a$, a set c and a $\Sigma_1(P_{\rho_0})$ -formula $A(x)$ such that $(\mathcal{H}, \kappa) \vdash_b^{a_0} \Gamma, \forall x \in c A(x)$ and $(\mathcal{H}, \kappa) \vdash_b^{a_0} \forall y \exists x \in c \neg A^{(y)}(x), \Gamma$, where for $A(x) \equiv (\exists z \in P_{\rho_0} \exists w B(x)) (B \in \Delta_0)$, $A^{(y)}(x) \equiv (\exists z \in P_{\rho_0} \cap y \exists w \in y B)$.

Lemma 5.12 (Tautology) If $\kappa(\Gamma \cup \{A\}) \subset \mathcal{H}$ then $(\mathcal{H}, \rho_0) \vdash_0^{2\text{rk}(A)} \Gamma, \neg A, A$.

Lemma 5.13 Let $\text{rk}(\forall x \in b \varphi[x, c]) \leq \rho_0 + m$ for an $m \geq 1$, and $\Theta_c = \{\neg \forall y (\forall x \in y \varphi[x, c] \rightarrow \varphi[y, c])\}$. Then for any operator \mathcal{H} , and any a, c , $(\mathcal{H}[\{c, a\}], \rho_0) \vdash_{\rho_0+m+1}^{\rho_0+2m+2+2\text{rk}_L(a)} \Theta_c, \forall x \in a \varphi[x, c]$.

Let

$$\begin{aligned} (\mathcal{H}, \rho_0) \vdash_c^{<\alpha} \Gamma & \Leftrightarrow \exists \beta < \alpha [(\mathcal{H}, \rho_0) \vdash_c^\beta \Gamma] \\ (\mathcal{H}, \rho_0) \vdash_{<c}^{<\alpha} \Gamma & \Leftrightarrow \exists d < c [(\mathcal{H}, \rho_0) \vdash_d^{<\alpha} \Gamma] \end{aligned}$$

Lemma 5.14 *Let A be an axiom in $\mathrm{T}(\omega_1)$ except Foundation axiom schema and Π_1 -Collection. Then $(\mathcal{H}, \rho_0) \vdash_0^{<\rho_0+\omega} A$ for any operator $\mathcal{H} = \mathcal{H}_\gamma$.*

Lemma 5.15 (Embedding)

If $\mathrm{T}(\omega_1) \vdash \Gamma[\vec{x}]$, there are $m, k < \omega$ such that for any $\vec{a} \subset L_{\rho_0}$, $(\mathcal{H}[\vec{a}], \rho_0) \vdash_{\rho_0+m}^{\rho_0 \cdot 2+k} \Gamma[\vec{a}]$ for any operator $\mathcal{H} = \mathcal{H}_\gamma$.

Proof.

By Lemma 5.13 we have $(\mathcal{H}, \rho_0) \vdash_{\rho_0+m+1}^{\rho_0 \cdot 2} \forall u, z (\forall y (\forall x \in y \varphi[x, z] \rightarrow \varphi[y, z]) \rightarrow \varphi[u, z])$ for some m . By Lemmata 5.12 and 5.14 it remains to consider instances

$$\forall u \in a \exists v \forall w \theta \rightarrow \exists c \forall u \in a \exists v \in c \forall w \theta$$

of Π_1 -Collection, where $\theta \equiv \theta(u, v, w)$ is a Δ_0 -formula in the language $\{\in\}$.

First by Lemma 5.14 with axioms (3) and (4) we have

$$(\mathcal{H}, \rho_0) \vdash_{\rho_0+1}^{\rho_0+\omega} \forall w \theta(u, v, w) \leftrightarrow \exists x \in P_{\rho_0} \tau(x, u, v)$$

where $\tau(x, u, v) \equiv [u, v \in L_x \wedge \forall w \in L_x \theta(u, v, w)]$. Hence

$$(\mathcal{H}, \rho_0) \vdash_{\rho_0+\omega}^{<\rho_0+\omega \cdot 2} \neg \forall u \in a \exists v \forall w \theta, \forall u \in a \exists x \in P_{\rho_0} \exists v \tau(x, u, v)$$

On the other hand we have by Lemma 5.12

$$(\mathcal{H}, \rho_0) \vdash_0^{<\rho_0+\omega} \neg \exists c \forall u \in a \exists x \in P_{\rho_0} \cap c \exists v \in c \tau, \exists c \forall u \in a \exists x \in P_{\rho_0} \cap c \exists v \in c \tau$$

Hence by the inference (*Ref*) for the $\Sigma_1(P_{\rho_0})$ -formula $\exists x \in P_{\rho_0} \exists v \tau(x, u, v)$, we obtain

$$(\mathcal{H}, \rho_0) \vdash_{\rho_0+\omega}^{<\rho_0+\omega \cdot 2} \neg \forall u \in a \exists v \forall w \theta, \exists c \forall u \in a \exists x \in P_{\rho_0} \cap c \exists v \in c \tau$$

Therefore $(\mathcal{H}, \rho_0) \vdash_{\rho_0+\omega}^{<\rho_0+\omega \cdot 2} \forall u \in a \exists v \forall w \theta \rightarrow \exists c \forall u \in a \exists v \in c \forall w \theta$. \square

In the following Lemma 5.16, note that $\mathrm{rk}(\exists x < \omega_1 \exists y < \omega_1 [\alpha < x \wedge P(x, y)]) = \omega_1 + 1$, and $\mathrm{rk}(\exists x < \rho_0 [\alpha < x \wedge P_{\rho_0}(x)]) = \rho_0$.

Lemma 5.16 (Predicative Cut-elimination)

1. If $(\mathcal{H}, \kappa) \vdash_{c+\omega^a}^b \Gamma \& [c, c + \omega^a[\cap\{\omega_1 + 1, \rho_0\}] = \emptyset \& a \in \mathcal{H} \Rightarrow (\mathcal{H}, \kappa) \vdash_c^{\varphi^{ab}} \Gamma$.
2. If $(\mathcal{H}_\gamma, \kappa) \vdash_{\omega_1+2}^b \Gamma \& \gamma \in \mathcal{H}_\gamma \Rightarrow (\mathcal{H}_{\gamma+b}, \kappa) \vdash_{\omega_1+1}^{\omega^b} \Gamma$.
3. If $(\mathcal{H}_\gamma, \kappa) \vdash_{\rho_0+1}^b \Gamma \& \gamma \in \mathcal{H}_\gamma \Rightarrow (\mathcal{H}_{\gamma+b}, \kappa) \vdash_{\rho_0}^{\omega^b} \Gamma$.

For a formula $\exists x \in d A(x)$ and ordinals $\lambda = \text{rk}_L(d), \alpha$, $(\exists x \in d A)^{(\exists \lambda \mid \alpha)}$ denotes the result of restricting the outermost existential quantifier $\exists x \in d$ to $\exists x \in L_\alpha$, $(\exists x \in d A)^{(\exists \lambda \mid \alpha)} \equiv (\exists x \in L_\alpha A)$.

Lemma 5.17 (Boundedness) *Let $\lambda \in \{\omega_1, \rho_0\}$, $C \equiv (\exists x \in d A) \in \Sigma^{\Sigma_2}(\lambda)$ and $C \notin \{\exists x < \omega_1 \exists y < \omega_1 [\alpha < x \wedge P(x, y)] : \alpha < \omega_1\} \cup \{\exists x < \rho_0 [\alpha < x \wedge P_{\rho_0}(x)] : \alpha < \rho_0\}$.*

1. $(\mathcal{H}, \lambda) \vdash_c^a \Lambda, C \& a \leq b \in \mathcal{H} \cap \lambda \Rightarrow (\mathcal{H}, \lambda) \vdash_c^a \Lambda, C^{(\exists \lambda \mid b)}$.
2. $(\mathcal{H}, \kappa) \vdash_c^a \Lambda, \neg C \& b \in \mathcal{H} \cap \lambda \Rightarrow (\mathcal{H}, \kappa) \vdash_c^a \Lambda, \neg(C^{(\exists \lambda \mid b)})$.

Lemma 5.18 (Boundedness for $\Sigma_1(P_{\rho_0})$)

Let C be a $\Sigma_1(P_{\rho_0})$ -sentence. Then $(\mathcal{H}, \rho_0) \vdash_c^a \Lambda, C \& a \leq b \in \mathcal{H} \cap \rho_0 \Rightarrow (\mathcal{H}, \rho_0) \vdash_c^a \Lambda, C^{(L_b)}$.

Proof. $C^{L_b} \equiv (\exists z \in P_{\rho_0} \cap L_b \exists w \in L_b B)$ when $C \equiv (\exists z \in P_{\rho_0} \exists w B)$ with a Δ_0 -formula B . The lemma is seen from (7). \square

5.4 Collapsing derivations

In this subsection derivations of $\Sigma^{\Sigma_2}(\omega_1)$ sentences are shown to be collapsed to derivations with heights and cut ranks $< \omega_1$.

Lemma 5.19 (Collapsing below ω_1)

Suppose $\gamma \in \mathcal{H}_\gamma[\Theta]$ with $\Theta \subset \mathcal{H}_\gamma(\Psi_{\omega_1}(\gamma))$, and $\Gamma \subset \Sigma^{\Sigma_2}(\omega_1)$. Then for $b = \Psi_{\omega_1}(\gamma + \omega^{\omega_1+a})$,

$$(\mathcal{H}_\gamma[\Theta], \omega_1) \vdash_{\omega_1+1}^a \Gamma \Rightarrow (\mathcal{H}_{\gamma+\omega^{\omega_1+a}+1}[\Theta], \omega_1) \vdash_b^b \Gamma.$$

Lemma 5.20 (Collapsing below ρ_0)

Suppose $\gamma \in \mathcal{H}_\gamma[\Theta]$ with $\Theta \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$, and $\Gamma \subset \Sigma^{\Sigma_2}(\rho_0) \cup \Sigma_1(P_{\rho_0})$. Then for $\hat{a} = \gamma + \omega^{\rho_0+a}$

$$(\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^a \Gamma \Rightarrow (\mathcal{H}_{\hat{a}+1}[\Theta], \rho_0) \vdash_{\Psi_{\rho_0}(\hat{a})}^{\Psi_{\rho_0}(\hat{a})} \Gamma.$$

Proof by induction on a , cf. Lemma 5.1.

First note that $\Psi_{\rho_0}(\hat{a}) \in \mathcal{H}_{\hat{a}+1}[\Theta]$ since $\hat{a} = \gamma + \omega^{\rho_0+a} \in \mathcal{H}_\gamma[\Theta] \subset \mathcal{H}_{\hat{a}+1}[\Theta]$ by the assumption, $\{\gamma, a\} \subset \mathcal{H}_\gamma[\Theta]$.

Assume $(\mathcal{H}_\gamma[\Theta_0], \rho_0) \vdash_{\rho_0}^{a_0} \Gamma_0$ with $\Theta_0 \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$. Then by $\gamma \leq \hat{a}$, we have $\hat{a}_0 \in \mathcal{H}_\gamma[\Theta_0] \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma)) \subset \mathcal{H}_{\hat{a}}(\Psi_{\rho_0}(\hat{a}))$. This yields that

$$a_0 < a \Rightarrow \Psi_{\rho_0}(\hat{a}_0) < \Psi_{\rho_0}(\hat{a})$$

Second observe that $\kappa(\Gamma) \subset \mathcal{H}_\gamma[\Theta] \subset \mathcal{H}_{\hat{a}+1}[\Theta]$ by $\gamma \leq \hat{a} + 1$.

Third we have

$$\kappa(\Gamma) \subset \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$$

When Γ is one of axioms (\mathbf{P}) and (\mathbf{P}_{ρ_0}) , there is nothing to show.

Consider the case when the last inference is a (Ref) .

$$\frac{(\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^{a_0} \Gamma, \forall x \in c A(x) \quad (\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^{a_0} \forall y \exists x \in c \neg A^{(y)}(x), \Gamma}{(\mathcal{H}_\gamma[\Theta], \rho_0) \vdash_{\rho_0}^a \Gamma} (Ref)$$

where $a_0 < a$ and $A(x) \equiv (\exists z \in P_{\rho_0} \exists w B(x))$ is a $\Sigma_1(P_{\rho_0})$ -formula with a Δ_0 -formula B .

For each $d \in c$ we have by Inversion

$$(\mathcal{H}_\gamma[\Theta \cup \{d\}], \rho_0) \vdash_{\rho_0}^{a_0} \Gamma, A(d)$$

where $c \in \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$. Hence $\text{rk}_L(d) < \text{rk}_L(c) \in \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma)) \cap \rho_0 \subset \Psi_{\rho_0}(\gamma)$, and $\text{rk}_L(d) < \Psi_{\rho_0}(\gamma)$. Therefore $d \in \mathcal{H}_\gamma(\Psi_{\rho_0}(\gamma))$. By IH we have for $\hat{a}_0 = \gamma + \omega^{\rho_0+a_0}$ and $\beta_0 = \Psi_{\rho_0}(\hat{a}_0) \in \mathcal{H}_{\hat{a}_0+1}[\Theta]$

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta \cup \{d\}], \rho_0) \vdash_{\beta_0}^{\beta_0} \Gamma, A(d)$$

Boundedness lemma 5.18 yields

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta \cup \{d\}], \rho_0) \vdash_{\beta_0}^{\beta_0} \Gamma, A^{(L_{\beta_0})}(d)$$

Since $d \in c$ is arbitrary, we obtain by (\wedge)

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta], \rho_0) \vdash_{\beta_0}^{\beta_0+1} \Gamma, \forall x \in c A^{(L_{\beta_0})}(x) \tag{8}$$

On the other hand we have by Inversion for $L_{\beta_0} \in \mathcal{H}_{\hat{a}_0+1}[\Theta]$

$$(\mathcal{H}_{\hat{a}_0+1}[\Theta], \rho_0) \vdash_{\rho_0}^{a_0} \exists x \in c \neg A^{(L_{\beta_0})}(x), \Gamma$$

Since $\exists x \in c \neg A^{(L_{\beta_0})}(x) \in \Sigma^{\Sigma_2}(\rho_0)$, IH yields for $\hat{a}_1 = \hat{a}_0 + 1 + \omega^{\rho_0+a_0} = \gamma + \omega^{\rho_0+a_0} \cdot 2$ and $\beta_1 = \Psi_{\rho_0}(\hat{a}_1)$

$$(\mathcal{H}_{\hat{a}_1+1}[\Theta], \rho_0) \vdash_{\beta_1}^{\beta_1} \exists x \in c \neg A^{(L_{\beta_0})}(x), \Gamma \tag{9}$$

We have $\text{rk}(\forall x \in c A^{(L_{\beta_0})}(x)) \in \text{Hull}(\kappa(\forall x \in c A^{(L_{\beta_0})}(x))) \cap \rho_0 \subset \mathcal{H}_{\hat{a}_0+1}[\Theta] \cap \rho_0 \subset \mathcal{H}_{\hat{a}_0+1}(\Psi_{\rho_0}(\gamma)) \cap \rho_0 \subset \Psi_{\rho_0}(\hat{a})$ by Proposition 5.10.5.

By a (cut) with (8) and (9) we obtain with $\Psi_{\rho_0}(\hat{a}) > \beta_1 > \beta_0$

$$(\mathcal{H}_{\hat{a}+1}[\Theta], \rho_0) \vdash_{\Psi_{\rho_0}(\hat{a})}^{\Psi_{\rho_0}(\hat{a})} \Gamma$$

Other case as seen as in [1]. □

6 Proof of Theorem 1.1

For a sentence $\exists x \in L_{\omega_1} \varphi$ with a Σ_2 -formula φ in the language $\{\in, \omega_1\}$, assume $T_1 \vdash \exists x \in L_{\omega_1} \varphi$. Then by Lemmata 3.2 and 5.15, pick an $m > 0$ such that the

fact $(\mathcal{H}_0, \rho_0) \vdash_{\rho_0+m}^{\rho_0 \cdot 2 + m} \exists x \in L_{\omega_1} \varphi$ is provable in $\text{FiX}^i(T_1)$. In what follows work in $\text{FiX}^i(T_1)$. Predicative Cut Elimination 5.16.1 and 5.16.3 yields

$$(\mathcal{H}_\gamma, \rho_0) \vdash_{\rho_0}^a \exists x \in L_{\omega_1} \varphi$$

for $\gamma = \omega_{m-1}(\rho_0 \cdot 2 + m)$ and $a = \omega_m(\rho_0 \cdot 2 + m)$. Then Collapsing below ρ_0 5.20 yields

$$(\mathcal{H}_{\omega_{m+1}(\rho_0 \cdot 2 + m) + 1}, \rho_0) \vdash_\beta^\beta \exists x \in L_{\omega_1} \varphi$$

for $\gamma + \omega^{\rho_0+a} = \omega_{m+1}(\rho_0 \cdot 2 + m)$ and $\beta = \Psi_{\rho_0}(\omega_{m+1}(\rho_0 \cdot 2 + m))$. Predicative Cut Elimination 5.16.1 and 5.16.2 yields

$$(\mathcal{H}_{\omega_{m+1}(\rho_2 \cdot 2 + m) + \varphi\beta\beta}, \rho_0) \vdash_{\omega_1+1}^{\varphi\beta\beta} \exists x \in L_{\omega_1} \varphi$$

for $\omega_1 + 2 + \omega^\beta = \beta$ and $\omega^{\varphi\beta\beta} = \varphi\beta\beta$. A fortiori,

$$(\mathcal{H}_{\omega_{m+1}(\rho_2 \cdot 2 + m) + \varphi\beta\beta}, \omega_1) \vdash_{\omega_1+1}^{\varphi\beta\beta} \exists x \in L_{\omega_1} \varphi$$

Then Collapsing below ω_1 5.19 yields

$$(\mathcal{H}_{\omega_{m+1}(\rho_2 \cdot 2 + m) + (\varphi\beta\beta \cdot 2) + 1}, \omega_1) \vdash_\delta^\delta \exists x \in L_{\omega_1} \varphi$$

for $\omega_{m+1}(\rho_2 \cdot 2 + m) + \varphi\beta\beta + \omega^{\omega_1 + \varphi\beta\beta} + 1 = \omega_{m+1}(\rho_2 \cdot 2 + m) + (\varphi\beta\beta \cdot 2) + 1$ and $\delta = \Psi_{\omega_1}(\omega_{m+1}(\rho_2 \cdot 2 + m) + (\varphi\beta\beta \cdot 2))$.

Boundedness 5.17.1 yields for $\delta < \Psi_{\omega_1}(\omega_n(\rho_0 + 1))$ with $n = m + 2$

$$(\mathcal{H}_{\omega_n(\rho_0+1)+1}, \omega_1) \vdash_\delta^\delta \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$$

We see then by transfinite induction up to the countable ordinal δ that inference rules in the controlled derivation of $\exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$ with cut rank $< \omega_1$ are (\vee) , (\wedge) , (cut) , and $(\mathbf{F}_{x \cup \{\omega_1\}})$, and since these inference rules are truth-preserving, we conclude again by transfinite induction up to δ that $\exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$ is true.

Since the whole proof is formalizable in $\text{FiX}^i(T_1)$, we conclude $\text{FiX}^i(T_1) \vdash \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$. Finally Theorem 5.2 yields $T_1 \vdash \exists x \in L_{\Psi_{\omega_1}(\omega_n(\rho_0+1))} \varphi$. This completes a proof of Theorem 1.1.

References

- [1] T. Arai, Lifting proof theory to the countable ordinals: Zermelo-Fraenkel's set theory, *Jour. Symb. Logic* **79** (2014), pp. 325-354.
- [2] T. Arai, Intuitionistic fixed point theories over set theories, *Arch. Math. Logic* **54** (2015), pp. 531-553.
- [3] W. Buchholz, A simplified version of local predicativity, P. H. G. Aczel, H. Simmons and S. S. Wainer(eds.), Proof Theory, Cambridge UP, 1992, pp. 115-147.